## Automata and Grammars

## SS 2018

## Assignment 3

Solutions are to be presented at the Seminary on Thursday, March 15, 2018.

## Problem 3.1 [Nerode Relation]

Determine all equivalence classes of the Nerode relations $R_{L} \subseteq \Sigma^{*} \times \Sigma^{*}$ for the following languages:
(a) $\Sigma=\{a, b\}$ and $L=\{a, a a b, a b b\}$,
(b) $\Sigma=\{a, b\}$ and $L=\left\{a^{m} b a^{n} \mid m, n \geq 1\right\}$,
(c) $\Sigma=\{a, b\}$ and $L=\left\{a^{n} b a^{n} \mid n \geq 0\right\}$,

In addition, for those language $L$, for which $R_{L}$ has finite index, construct the equivalence class automaton for $L$.

Problem 3.2 [Myhill Nerode Theorem]
Use the Myhill Nerode Theorem to check which of the following languages are accepted by DFAs:
(a) $L_{a}=\left\{a^{n} b^{n} \mid n \geq 1\right\}$,
(b) $L_{b}=\left\{w w \mid w \in\{a, b\}^{*}\right\}$,
(c) $L_{c}=\left\{a^{2 n} \mid n \geq 0\right\}$,
(d) $L_{d}=\left\{a^{2^{n}} \mid n \geq 0\right\}$.

Problem 3.3 [DFAs]
Two states $p, q$ of a DFA $A=\left(Q, \Sigma, \delta, q_{0}, F\right)$ are called equivalent if, for all $w \in \Sigma^{*}, \delta(p, w) \in$ $F$ iff $\delta(q, w) \in F$. Find all pairs of equivalent states in the following four DFAs:
(a) $A_{1}=\left(\left\{q_{0}, q_{1}, \ldots, q_{7}\right\},\{a, b\}, \delta_{1}, q_{0},\left\{q_{0}, q_{4}, q_{6}\right\}\right)$, where $\delta_{1}$ is described through the following table:

| $\delta_{1}$ | $q_{0}$ | $q_{1}$ | $q_{2}$ | $q_{3}$ | $q_{4}$ | $q_{5}$ | $q_{6}$ | $q_{7}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $a$ | $q_{0}$ | $q_{1}$ | $q_{2}$ | $q_{3}$ | $q_{6}$ | $q_{5}$ | $q_{4}$ | $q_{7}$ |
| $b$ | $q_{5}$ | $q_{3}$ | $q_{7}$ | $q_{2}$ | $q_{1}$ | $q_{1}$ | $q_{2}$ | $q_{0}$ |

(b) $A_{2}=\left(\{A, B, C, D, E, F\},\{a, b\}, \delta_{2}, F,\{F\}\right)$, where $\delta_{2}$ is described by the following table:

| $\delta_{2}$ | $A$ | $B$ | $C$ | $D$ | $E$ | $F$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $a$ | $A$ | $B$ | $C$ | $D$ | $E$ | $F$ |
| $b$ | $F$ | $A$ | $D$ | $B$ | $C$ | $E$ |

(c) $A_{3}=\left(\left\{q_{1}, q_{2}, \ldots, q_{9}\right\},\{a, b\}, \delta_{1}, q_{1},\left\{q_{3}, q_{5}, q_{6}\right\}\right)$, where $\delta_{3}$ is described through the following table:

| $\delta_{3}$ | $q_{1}$ | $q_{2}$ | $q_{3}$ | $q_{4}$ | $q_{5}$ | $q_{6}$ | $q_{7}$ | $q_{8}$ | $q_{9}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $a$ | $q_{2}$ | $q_{2}$ | $q_{3}$ | $q_{2}$ | $q_{6}$ | $q_{6}$ | $q_{7}$ | $q_{2}$ | $q_{9}$ |
| $b$ | $q_{3}$ | $q_{4}$ | $q_{5}$ | $q_{7}$ | $q_{3}$ | $q_{6}$ | $q_{4}$ | $q_{3}$ | $q_{4}$ |

(d) $A_{4}=\left(\{A, B, C, D, E, F, G, H\},\{a, b\}, \delta_{4}, G,\{G\}\right)$, where $\delta_{4}$ is described by the following table:

| $\delta_{4}$ | $A$ | $B$ | $C$ | $D$ | $E$ | $F$ | $G$ | $H$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $a$ | $H$ | $B$ | $E$ | $D$ | $C$ | $F$ | $G$ | $A$ |
| $b$ | $G$ | $A$ | $D$ | $B$ | $D$ | $E$ | $F$ | $G$ |

## Problem 3.4 [Minimal DFA]

Use the algorithm "minimal automaton" to construct the minimal DFAs that are equivalent to the following DFAs:
(a) $B_{1}=\left(\left\{q_{1}, q_{2}, \ldots, q_{9}\right\},\{a, b\}, \delta_{1}, q_{1},\left\{q_{3}, q_{5}, q_{6}\right\}\right)$, where $\delta_{1}$ is described through the following table:

| $\delta_{1}$ | $q_{1}$ | $q_{2}$ | $q_{3}$ | $q_{4}$ | $q_{5}$ | $q_{6}$ | $q_{7}$ | $q_{8}$ | $q_{9}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $a$ | $q_{2}$ | $q_{2}$ | $q_{3}$ | $q_{2}$ | $q_{6}$ | $q_{6}$ | $q_{7}$ | $q_{2}$ | $q_{9}$ |
| $b$ | $q_{3}$ | $q_{4}$ | $q_{5}$ | $q_{7}$ | $q_{3}$ | $q_{6}$ | $q_{4}$ | $q_{3}$ | $q_{4}$ |

(b) $B_{2}=\left(\{A, B, C, D, E, F, G, H\},\{a, b\}, \delta_{2}, G,\{G\}\right)$, where $\delta_{2}$ is described by the following table:

| $\delta_{2}$ | $A$ | $B$ | $C$ | $D$ | $E$ | $F$ | $G$ | $H$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $a$ | $H$ | $B$ | $E$ | $D$ | $C$ | $F$ | $G$ | $A$ |
| $b$ | $G$ | $A$ | $D$ | $B$ | $D$ | $E$ | $F$ | $G$ |

(c) $B_{3}=\left(\left\{q_{0}, q_{1}, \ldots, q_{5}\right\},\{a, b\}, \delta_{3}, q_{0},\left\{q_{0}\right\}\right)$, where $\delta_{3}$ is described through the following table:

| $\delta_{3}$ | $q_{0}$ | $q_{1}$ | $q_{2}$ | $q_{3}$ | $q_{4}$ | $q_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $a$ | $q_{1}$ | $q_{3}$ | $q_{4}$ | $q_{0}$ | $q_{2}$ | $q_{0}$ |
| $b$ | $q_{2}$ | $q_{0}$ | $q_{5}$ | $q_{2}$ | $q_{5}$ | $q_{3}$ |

